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(71) Applicant: **Hewlett-Packard Company**
Mail Stop 20 B-O, 3000 Hanover Street
Palo Alto, California 94304(US)

(72) Inventor: **Phaal, Peter**
25, The Valls Bradley Stoke
Bristol(GB)

(74) Representative: **Squibbs, Robert Francis et al**
Hepworth, Lawrence, Bryer & Bizley, Lewins House, Lewins Mead
Bristol BS1 2NN(GB)

(54) Method of ascertaining topology features of a network.

(57) A method is provided for ascertaining topology features of a network of the type comprising a plurality of sub-networks, spanning devices (14) interconnecting the sub-networks, and stations (11) operative to source and/or sink traffic to/from the sub-networks, the traffic being in the form of discrete message packets each including source and destination information. Traffic on each of the sub-networks is monitored by a respective monitoring device (12) and a partial traffic matrix is built up for each sub-network. These partial traffic matrices are then analysed by a central processing station 13 to determine, for example, the association of stations (11) with sub-networks, and the interconnection of pairings of sub-networks by spanning devices (14).

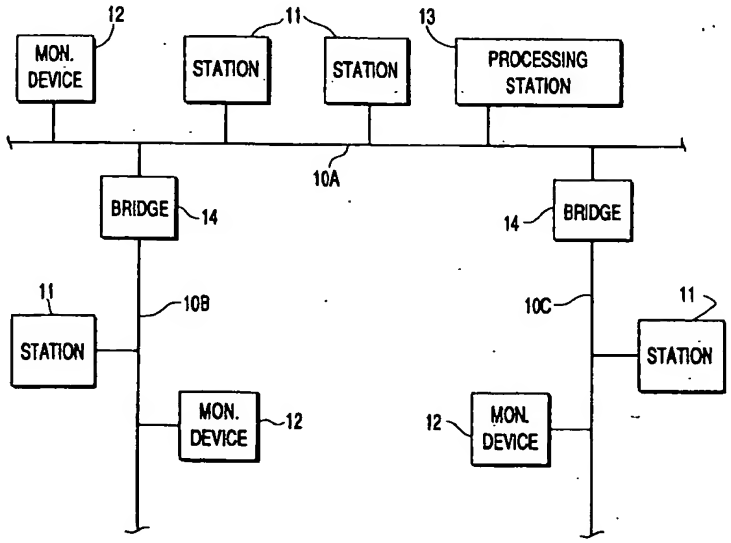

FIG 1

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The present invention relates to a method of ascertaining topology features of a network of the type comprising a plurality of sub-networks, spanning devices interconnecting the sub-networks, and stations operative to source and/or sink traffic to/from the sub-networks, the traffic being in the form of discrete message packets each including source and destination information.

Networks of the aforesaid type are well known, their spanning devices serving to isolate local traffic (that is, traffic sourcing from, and destined for, stations on the same sub-network) to the sub-network concerned, at least to some degree. Examples of such networks are:

(a) - bridged Ethernets where stations, identified by Ethernet addresses, are connected to sub-networks (also termed logical segments) which are interconnected by spanning devices in the form of bridges operating at level 2 of the seven-layer OSI Reference Model; and

(b) - Internet networks where stations identified by "ip" addresses are connected to "ip" sub-networks which are interconnected by spanning devices in the form of routers or gateways operating at level 3 of the seven-layer OSI Reference Model.

Knowledge of topological features of such networks (for example, the connection arrangement of stations to sub-networks and the interconnection of sub-networks by spanning devices) is of importance in monitoring and optimising the performance of the network and in planning for its expansion to cope with increased demand. However, such knowledge is frequently difficult to ascertain, particularly with large networks. For example, keeping up-to-date plans of the network by recording every change made, is not only time-consuming but is, in reality, a virtually impossible task to carry through with complete accuracy; in addition, such plans can only indicate the intended connection state of the network and cannot take account of the failure of network elements such as the simple disconnection of a station from its sub-network. Physical inspection of a network to ascertain its topology is also very difficult since much of the network may be hidden beneath flooring or located at remote sites.

In the case of a network made up of a single sub-network, it is possible to ascertain which stations are connected to the subnetwork simply by monitoring the traffic on the sub-network; all stations sourcing traffic are obviously connected to the sub-network as there is no other source of traffic outside the sub-network. However, where the network comprises several sub-networks interconnected by spanning devices, it is no longer possible to make such simple deductions as traffic appearing on any one particular sub-network might have originated on another sub-network.

It is an object of the present invention to provide a relatively simple method for ascertaining topology features of a network that includes a plurality of sub-networks.

According to the present invention, there is provided a method of ascertaining topology features of a network of the aforesaid type, characterised in that said method comprises the steps of:

- monitoring the traffic on at least some of said sub-networks to derive data concerning at least the amounts and origins of the traffic so monitored, and
- processing the said data to determine selected topology features of the network.

The monitoring step utilises a monitoring device on each sub-network concerned. Such a device may be a stand-alone device or part of a spanning device or station connected to the sub-network. Each monitoring device is preferably operative to collect both source and destination information in respect of traffic on the sub-network and this information is then used to generate a partial traffic matrix, that is, a correlation between source stations and destination stations for the traffic seen on the sub-network concerned. Preferably, therefore, the data passed by the monitoring step to the processing step will be in the form of partial traffic matrices for all the relevant sub-networks. However, it should be understood that the processing step, depending on what topological features are to be ascertained, may not require all the data contained in such partial traffic matrices and the present invention should therefore not be taken as limited to the provision of the necessary data to the processing step in the form of partial traffic matrices.

The source information in each message packet will generally only provide an indication of the originating station, although in certain networks, sub-network identification is also carried in the source information. In the former case, the method of the present invention can be used to ascertain to which sub-network a selected station is connected; this is achieved by having said monitoring step determine the amount of traffic originating from said selected station that is carried by each said sub-network, the said processing step then determining the sub-network carrying the most traffic from said selected station, this sub-network being taken as the one to which said selected station is connected. This process can then be repeated for each station to build up a complete picture of the connection of stations to sub-networks.

The method of the present invention can also be used to ascertain whether a first sub-network is directly connected by a spanning device to a second sub-network; this is done by comparing, for some or all of the components of traffic originating

on the first sub-network, the amounts of that traffic seen on the second sub-network with the amounts of that traffic seen on all the other sub-networks (that is, other than the first and second sub-networks). In particular, three different methods for testing for a spanning device are provided.

In the first of the methods for testing for a spanning device between first and second sub-networks:

- said method includes the step of providing association information enabling the said source information included in a message packet to be used to identify a message packet as having originated from said first sub-network;
- said monitoring step involves monitoring the traffic on said second sub-network and on each sub-network of the group of sub-networks made up of all said sub-networks other than said first and second sub-networks, whereby to derive, with the use of said association information, data on the amounts of traffic sourcing from said first sub-network that are carried by said second sub-network and the sub-networks of said group of sub-networks; and
- said processing step is operative to use said data to compare the said amounts of traffic sourcing from said first sub-network, that are respectively carried by said second sub-network and by the sub-networks of said group taken together, the processing step determining that a spanning device directly connects the first sub-network to the second when the larger of these amounts of traffic is carried by said second sub-network.

In the second of the methods for testing for a spanning device between first and second sub-networks:

- said method includes the step of providing association information enabling:
  (a) the said source information included in a message packet to be used to identify a message packet as having originated from said first sub-network; and
  (b) the said destination information included in a message packet to be used to identify a message packet as having a destination other than said second sub-network;
- said monitoring step involves monitoring the traffic on said second sub-network and on each sub-network of the group of sub-networks made up of all said sub-networks other than said first and second sub-networks, whereby to derive, with the use of said association information, data on the amounts of traffic sourcing from said first sub-network and having a destination other than on said second sub-network, that are carried by said second sub-network and the sub-networks of said group of sub-networks; and
- said processing step is operative to use said data to compare the said amounts of traffic sourcing from said first sub-network and having a destination other than on said second sub-network, that are respectively carried by said second sub-network and by the sub-networks of said group taken together, the processing step determining that a spanning device directly connects the first sub-network to the second when the larger of these amounts of traffic is carried by said second sub-network.

The third one of the methods for testing for a spanning device between first and second sub-networks is applicable where the spanning devices used in the network do not act perfectly to block messages packets with a destination on a connected sub-network from passing through the spanning device from that sub-network to another. In this third method for testing for a spanning device:

- said method includes the step of providing association information enabling:
  (a) the said source information included in a message packet to be used to identify a message packet as having originated from said first sub-network; and
  (b) the said destination information included in a message packet to be used to identify a message packet as having a destination on said first sub-network;
- said monitoring step involves monitoring the traffic on said second sub-network and on each sub-network of the group of sub-networks made up of all said sub-networks other than said first and second sub-networks, whereby to derive, with the use of said association information, data on the amounts of traffic sourcing from said first sub-network and having a destination on said first sub-network, that are carried by said second sub-network and the sub-networks of said group of sub-networks; and
- said processing step is operative to use said data to compare the said amounts of traffic sourcing from said first sub-network and having a destination on said first sub-network, that are respectively carried by said second sub-network and by the sub-networks of said group taken together, the processing step determining that a spanning device directly connects the first sub-network to the second when the larger of these amounts of traffic is carried by said second sub-network.

All three of the foregoing methods of testing for a spanning device are conservative in operation in

that they may not find a spanning device which is present, but they are also safe in that they will not find a spanning device which is not present.

The association information required by all three of the methods for testing for a spanning device, associates stations with sub-networks. Where the source information in each message packet inherently provides an indication of the sourcing sub-network, then the association information is simply that required to interpret the source information correctly according to some predetermined scheme; in other cases, the ad hoc association between stations and sub-networks must first be determined (for example, in accordance with the above-recited method of the invention) and recorded as said association information.

In all three of the methods for testing for a spanning device, the derivation in the processing step of the relevant amount of traffic sourcing from said first sub-network that is carried by the sub-networks of said group taken together, can be effected in several different ways. For example, the derivation can be effected by adding together the corresponding traffic amounts for each sub-network of said group without regard to any duplication that may occur as a result of the same message packet being monitored on more than one sub-network of the group. Alternatively the derivation can be effected by:

- determining for each sub-network of said group, the components of the sub-network traffic sourcing from said first sub-network, that are destined for that sub-network and for said second sub-network, and
- adding together the sub-network traffic components so determined for all the sub-networks of said group.

Any of the three aforesaid methods of testing for a spanning device can be effected in relation to traffic sourcing from said first sub-network either without regard to originating station on the sub-network (that is, by taking all stations together), or by taking each station separately and applying the test to traffic sourcing from that station on the first sub-network.

Furthermore, where bi-directional spanning devices are used (that is, they permit traffic flow in both directions between the sub-networks they interconnect) each of the tests for a spanning device between first and second sub-networks can be carried out not only for traffic sourcing from the first network but also for traffic sourcing from the other said sub-network.

It is, of course, possible to use all three methods of testing for a spanning device, in succession to test for a spanning device between the same pair of sub-networks (testing being stopped when and if a positive result is obtained). By applying the aforesaid methods of testing for a spanning device to each pairing of sub-networks of the network or the portion thereof under consideration, it is possible to determine which of the sub-networks are interconnected by spanning devices and thereby to build up a full picture of the network topology.

In a preferred implementation of the invention, traffic information on each sub-network is collected locally by a sampling monitoring device which transmits sampled data back to a central processing station. The processing station derives partial traffic matrices for all the monitored sub-networks and then processes the data contained in the traffic matrices in accordance with the methods of the invention. In carrying out this processing, the processing station effects comparisons between traffic flows by using hypothesis testing in view of the sampled nature of the data provided to it.

A suitable sampling monitoring device for use in implementing the method of the present invention is described and claimed in our co-pending European Patent Application No.90310699.5 filed 28th September 1990.

The partial traffic matrices used by the processing station for deriving topology features of the network can, of course, be obtained in other ways such as, for example, by monitoring devices which derive and store these partial traffic matrices themselves, the partial traffic matrices then being collected and passed to the processing station.

A method, according to the invention, of ascertaining topology features of a network will now be described by way of non-limiting example with reference to the accompanying diagrammatic drawings, in which:

Figure 1 is an overall diagram of a network to which a processing station and a number of sampling monitoring devices have been connected to form a network monitoring system implementing the method of the invention;

Figure 2 is a diagram illustrating the general form of a data packet transmitted over the Figure 1 network;

Figure 3 is a block diagram of a sampling monitoring device of Figure 1;

Figure 4 is a flow chart illustrating the main interrupt service routine run by a controlling microprocessor of the Figure 3 device;

Figure 5 is a diagram illustrating the main data structures utilised by the processing station of Figure 1 in processing data from the sampling monitoring devices;

Figure 6 shows an example partial traffic matrix for a sub-network of the

Figure 1 network;

Figure 7 is a flow chart illustrating a main program run by the processing station;

Figure 8 is a flow chart illustrating an interrupt service routine which is run by the processing station, upon receipt of data from a sampling monitoring device, in order to build up partial traffic matrices;

Figure 9 is a flow chart illustrating a "station-allocation" routine called by the Figure 7 program to determine the association between stations and sub-networks of the Figure 1 network;

Figure 10A is a diagram illustrating possible traffic flows in the network between two sub-networks Y and X when they are connected by a bridge;

Figure 10B is a diagram similar to Figure 10A but for the case when no bridge inter-connects sub-networks Y and X.

Figure 11 is a diagram illustrating possible leakage traffic flows from sub-network Y when no bridge interconnects sub-networks Y and X.

Figure 12 is a flow chart illustrating a "tests for bridges" routine called by the Figure 7 program to determine which pairs of sub-networks are inter-connected by bridges; and

Figure 13 is a flow chart illustrating a "first test for a bridge" sub-routine called by the Figure 12 routine.

Figure 1 illustrates a typical local area network in which a plurality of stations 11, 12, and 13 are interconnected via cable segments 10A, 10B, and 10C. The network is divided into three sub-networks by bridges (spanning devices) 14 that connect respective ones of the cable segments 10B, 10C to the cable segment 10A. As is well known in the art, the bridges serve to filter traffic passing between the network segments, such that messages originating from a particular segment and destined for a station on the same segment (local traffic) are not passed through the bridge or bridges 14 to the other segments whereas messages originating in one segment and intended for another one (non-local traffic) are allowed across the bridge. The operation for such bridges is generally imperfect and some local traffic will usually 'leak' through the bridges.

In the illustrated local area network, messages between the stations 11, 12 and 13 are transmitted in the form of packets that are broadcast over the network. Typically a packet will have the form illustrated in Figure 2 with a packet header 15 containing a source address (the address of the station sending the packet) and a destination address (the address of the station intended to receive the packet), and an information field 16 containing the data to be passed to the receiving station and normally including error checking codes. Depending on the particular packet format being used, other fields may also be present; thus, for example, there may be a CRC (cycle redundancy check) field covering both the packet header and information field.

The Figure 1 network may, for example, be an Ethernet network well known to persons skilled in the art.

The network of Figure 1 is arranged to be monitored by a network monitoring system comprising a plurality of monitoring devices (stations 12) and a central processing station 13. Each of the monitoring devices is associated with a respective one of the sub-networks of the network. As will become clear below, each monitoring device is operative to randomly sample the packets on its associated sub-network and transmit data on the sampled packets back to the processing station 13 for processing and analysis.

The form of each monitoring device is illustrated in Figure 3. The device comprises a network interface 20, a microprocessor 21, and ROM (non-volatile, pre-programmed memory) and RAM (re-writable memory) units 22 and 23. These units 20 to 23 are all interconnected via address, data and control buses 27, 28 and 29. The network interface 20 is operative to carry out all the low level functions necessary to interface the monitoring device of Figure 3 to the network cable 10 and to pass received packets to a receive queue, in the form of a FIFO (First In First Out) buffer 25 in RAM 23. The network interface is further operative to transmit packets held in a transmit queue, formed by a FIFO buffer 26, in RAM 23. The network interface 20 thus constitutes packet receive means and packet transmit means for the monitoring device. In the present example, the network interface 20 is arranged to receive all packets regardless of their destination address contained in the packet header. Furthermore, the network interface 20 is operative to pass only the header portion 30 of each received packet to the receive FIFO buffer 25.

The network interface 20 is arranged to operate in coordination with the microprocessor controller 21 and, in particular, informs the microprocessor 21 each time a packet header is inserted into the receive FIFO buffer 25, by means of a suitable interrupt control signal.

The network interface 20 also contains various

counters 24 which hold a number of counts including the total number of packets received, the number of packets received which according to their CRC field are in error, the number of packets received below the minimum accepted length (RUNT packets), and the number of packets received above the maximum accepted length (JABBER).

Implementations of the network interface 20 for particular network protocols are well known in the art. Thus, for example, for an Ethernet network, the network interface 20 may be constituted by Intel Corporation chips 82502, 82501, and 82586; in this case an appropriate microprocessor constituting the microprocessor 21 is the Intel processor 80186.

The ROM 22 holds the programs run by the microprocessor 21 and also a table of random count values predetermined according to an exponential distribution.

The processor 21 is operative to run a background program in which it does nothing (ie an idling program). The main working program for the processor 21 is an interrupt service routine which is called each time the network interface 20 generates a processor interrupt to tell the processor that it has stored a new packet header in the receive FIFO 25. The interrupt service routine, which will be described in more detail below, operates to randomly select a received packet header and form it into a collected-data packet together with the current count values of the counters 24; the random selection of received packet headers is effected by utilizing the predetermined random counts stored in ROM 22. The collected-data packet so formed is put into the transmit queue FIFO 26 and, in due course, is transmitted by the network interface 20 back to the processing station 13. The header of each collected-data packet contains as its source address the address of the monitoring device concerned while the destination address is that of the processing station (alternatively, a multicast address can be used to which the processing station is set to listen).

A more detailed description of the operation of the monitoring device will now be given with reference to Figure 4 which is a flow chart of the interrupt service routine run by the microprocessor 21. The microprocessor 21 will be taken to be in a state in which it is running its background (idling) program and in which it has one of the random count values held in an internal register (the fetching of the first count value upon switch-on of the monitoring device would be part of an initialization routine). It will also be assumed that the receive and transmit FIFO buffers 25 and 26 are empty.

On receiving a packet over the network cable 10, the network interface 20 passes the packet header to the receive FIFO buffer 25, updates its counters 24 and generates an interrupt signal for the microprocessor 21. On receipt of this interrupt, the microprocessor 21 executes the interrupt service routine illustrated in Figure 4. The first step 40 of this routine carries out the normal house-keeping tasks associated with such routines including saving the volatile environment parameters of the background program and masking further interrupts.

Next, the microprocessor decrements the random count value held in its internal register (step 41) and then checks the remaining value to see if this has been reduced to zero (step 42).

If the count value is still greater than zero, the microprocessor 21 discards the head entry in the receive FIFO buffer 25 (step 43).

Thereafter, the microprocessor must check the receive FIFO buffer 25 to see if any further packet headers have been entered into the buffer by the network interface 20 during the preceding steps of the interrupt service routine (step 44). Generally this will not be the case and the microprocessor will then exit its interrupt service routine and restore its background environment and unmask its interrupts (step 45). However, in the event that the receive FIFO buffer 25 contains a further packet header, the interrupt service routine will pass from step 44 back to step 41.

If during the test (step 42) carried out on the count value held in its internal register, the microprocessor 21 finds that this count value has been reduced to zero, the interrupt service routine will proceed to generate a collected-data packet 31 in respect of the packet header at the top of the receive FIFO buffer 25 (step 46). This collected-data packet 31 is assembled in the transmit FIFO buffer 26 from the received packet header 30, the count values from the counter 24, the address of the monitoring device (source address for the collected-data packet) and the address of the processing station (destination address for the collected-data packet header). After the collected-data packet has been assembled, the microprocessor 21 flags the network interface 20 to indicate that there is a packet ready for transmission. (The network interface 20 will transmit the packet as and when it is able and cancel the flag set by the microprocessor 21 once this has been done).

After completion of step 46 of the interrupt service routine, the microprocessor fetches a new random count from ROM 22 and stores this new random count in its internal register (step 47). The microprocessor then proceeds to step 44 and running of the interrupt service routine proceeds as previously described.

The size of the receive and transmit FIFO buffers 25 and 26 can be quite small, for example, sufficient to hold only two or three entries. This is

possible with respect to the receive buffer 25 because in general the interval between packets received by the network interface 20 will be sufficient for the microprocessor 21 to run its interrupt service routine and clear the top entry from the receive buffer; in any event, the occasional overflowing of the receive buffer 25 is not of major consequence since the missing out of a packet will generally have minimal effect on the statistical measurements being conducted by the network monitoring system. This equally applies to the transmit buffer 26 where an overflow is even less likely to occur as its entries are only in respect of the randomly selected ones of the received packets.

The above-described implementation of the monitoring device does mean that the count values included in a collected-data packet from the counter 24 may not be the count values current at the time that the relevant packet was actually received by the network interface (this is because of the possible delay in actually processing the packet header). However, again, any discrepancy in this respect will be minor and will have minimal effect on the validity of the statistically determined results produced by the network monitoring system. Of course, it would be possible to design circuitry which associated the count values present in counters 24 with the header of each received packet; however, the added circuit complexity needed to do this is generally not justified.

The data structures used to implement the receive and transmit FIFO buffers 25 and 26 in RAM 23 will be apparent to a person skilled in the art and will therefore not be described herein. Furthermore, it will be appreciated that although in the Figure 3 embodiment the random selection of incoming packets has been effected by storing predetermined random numbers in ROM 22, these random numbers could alternatively be generated as and when required by the processor 21 (although this is not preferred as it places extra processor requirements on the microprocessor). Typically, the random numbers are such as to give an average skip between selected packets of ninety nine; other values may be more appropriate depending on traffic density, sampling period and acceptable statistical error level. The random selection of packets could be effected on a time basis rather than on the number of packets received.

The collected-data packets sent out by the monitoring devices 12 over the network are all received by the processing station 13 which stores these packets and carries out subsequent processing and analysis.

The processing station 13 is, for example, constituted by a standard workstation interfacing with the network through a network interface (not shown) similar to the interface 20 of the Figure 3 monitoring device 12. Such a workstation will be provided in standard manner with RAM memory for storing working data and program segments, ROM memory for permanent storage of programs, a processor for processing data held in the RAM memory in accordance with the programs, and various input/output devices; none of these elements are illustrated or described herein as they are all standard and well known to persons skilled in the art.

The processing station 13 carries out three main tasks, namely:

(1) - generation of traffic matrices for each sub-network on the basis of the collected-data packets received;

(2) - the association of stations 11 with the various sub-networks using the sub-network traffic matrices; and

(3) - testing for the presence of bridge between all pairs of sub-networks by using the sub-network traffic matrices.

The main data structures employed by the processing station 13 in carrying out these tasks are illustrated in Figure 5, the data structures being created and maintained in the RAM memory of the station 13. The main data structures are:

Input-queue 50

- this is a queue of collected-data packets required by the network interface of the station, each packet being temporarily held in the queue pending processing;

Sub-Network List 51

- this is a list of all the known sub-networks of the network with each sub-network having a respective entry comprising a field storing the sub-network identity SN-ID, a first pointer TM-POINTER, and a second pointer S-POINTER;

Station List 52

- this is a list of all the known stations 11 of the network with each station having a respective entry comprising a field storing the station identity S-ID, and a pointer NEXT-S POINTER. The first station to be associated with any particular sub-network is associated with that sub-network by selling the S-POINTER of the corresponding sub-network entry in the sub-network list 51, to point to the appropriate station entry in the station list 52. The association of further stations with the same sub-network is achieved by using the NEXT-S POINTER of the last preceding station associated with the sub-network to point to the entry of the next station to be associated with the sub-network, thereby building up a linked list of stations.

Sub-Network Traffic Matrix 53

- This is an array formed for each sub-network to hold the partial traffic matrix data for the sub-network. The traffic matrix array relevant to a

sub-network is pointed to by the pointer TM-POINTER of the corresponding entry in the sub-network list 51. Figure 6 illustrates a typical partial traffic matrix giving for each source station/destination station pair, the number of packets carried by the sub-network concerned in a given interval of time (the different stations 11 are here designated 11A, 11B, 11C ..... 11N).

Bridge List 54

- This is a list of all the known bridges of the network with each bridge having a respective entry comprising a field storing the bridge identity B-ID, and a pair of fields for storing the identities (SN-ID) of the sub-networks connected by the bridge.

Having outlined the main data structures used by the processing station 13, a description will now be given as to how the station carries out its tasks using these structures.

The main program run by the processing station 13 is illustrated in Figure 7. Upon start up of the program, a predetermined time period, e.g. 1 hour, is timed (steps 60,61) during which the main program is idle but collected-data packets are received by the processing station and processed by an interrupt service routine to construct traffic matrices for the sub-networks of the network. After time-out of the predetermined time interval, the main program carries out statistical processing of the traffic matrices (step 62) before proceeding to ascertain topological features of the network, first of all by associating stations with sub-networks using a "station allocation" routine (step 63), and then by ascertaining where bridges are present in the network using the "tests for bridges" routine (step 64); finally, the main program terminates.

Figure 8 illustrates the interrupt service routine used to derive sub-network traffic matrices from the received collected-data packets. The interrupt service routine is called in response to the generation of an interrupt by the network interface of the station 13 following the receipt thereby of a collected-data packet and the storage of that packet in the input queue data structure 50. Upon start-up of the interrupt service routine, the normal housekeeping tasks of storing the volatile environment of the main program and masking interrupts is effected (step 65). Next, the header of the collected-data packet at the top of the input queue is examined (step 66). The source address contained in the header identifies the monitoring device that generated the collected-data packet; indeed, on the assumption that one monitoring device is associated with each sub-network of the network, the address of the monitoring device also serves to identify the associated sub-network. The monitoring device addresses are therefore used either directly or indirectly as the sub-network identities SN-ID. By examining the source address contained in the header of the collected-data packet under examination, it is therefore possible to tell, by checking against the sub-network list 51, whether the collected-data packet comes from a known sub-network or whether the packet indicates the presence of a previously unknown sub-network (step 67). If a new sub-network is identified, then a new entry is made in the sub-network list 51 and a new sub network traffic matrix array 53 is created (step 68); the TM-POINTER of the sub-network entry is set to point to the new traffic matrix array whilst the S-POINTER is set to null.

Thereafter, the interrupt service routine goes on to examine the contents of the data field of the collected-data package under consideration (step 69). A check is made of the source station and destination station addresses held in the data field to ascertain if either of these addresses indicates the presence of a previously unknown station, that is, a station not contained in the station list 52 (step 70). If a new station is identified, then the station list 52 is updated (with its NEXT-S POINTER set to null) and the existing sub-network traffic matrices 53 are extended to include the new station (step 71).

Once the foregoing preliminary steps have been effected, the interrupt service routine updates the traffic matrix 53 for the relevant sub-network (as identified by the source address in the header of the collected-data packet) by incrementing the entry for the source station/destination station pair concerned (as indicated by the addresses held in the data field of the colleted-data packet) - see step 72.

Because of the sampled nature of the information provided by the monitoring devices 12, the sub-network traffic matrices built directly using the raw data from the monitoring devices 12, is preferably subject to statistical processing. This processing is effected in step 62 of the main program at the time-out of the sampling period. However, in order to carry out this processing, it is necessary to record the total number of packets carried by each sub-network during the sampling period. As explained above, the total current packet count for a sub-network is included in the information field of each collected-data packet sent to the processing station 13. In order to ascertain the total number of packets carried on a sub-network during the sampling period, step 73 of the interrupt service routine is operative to record both the current packet count for the first collected-data packet received in respect of a sub-network and also the current packet count of the most recent collected-data packet received during the sampling period in respect of that sub-network. These count values can be re-

corded for each network by appropriate fields (not shown) in the corresponding sub-network entry in the sub-network list 51. At the end of the sampling period, subtraction of the count values will give the total number of packets carried by the sub-network during the sampling period.

Finally, the input queue 50 is checked to ensure that no new collected-data packets have been received whilst the interrupt service routine has been running (step 74); assuming this is the case, the interrupt service routine terminates (step 75) by restoring the main program environment parameters and unmasking the interrupts. In the event that the input queue is not empty, steps 66 to 73 are repeated.

Following the termination of the sampling period, the main program calls the routine for effecting statistical processing of the sub-network traffic matrices (step 62 of Figure 7). This processing involves deriving for each element of each sub-network traffic matrix, mean and variance values for an assumed normal distribution, these values being stored back into the relevant traffic matrix array. In order to calculate these mean and variance values, it is necessary to know not only the basic aggregated number of relevant samples "o_c" constituted by each traffic matrix element (that is, the number of samples of a class"C") , but also the total number "o" of samples involved in the generation of the relevant sub-network traffic matrix, and the total number of packets "f" carried by the sub-network during the sampling period. The total number "o" of samples can be obtained either by summing all elements of the relevant sub-network traffic matrix or by keeping a running total during the sampling period. As already discussed, the total number of packets "f" is obtained by subtracting an initial packet count value from a final count value, both these values being stored in the sub-network list 51. The values "o_c", "o" and "f" are then used in accordance with the following

$$\text{estimate of mean} = fo_c/o \quad (1)$$

$$\text{estimate of variance} = f^2(o-o_c)o_c/o^3 \quad (2)$$

Following this statistical processing of the sub-network traffic matrix elements, the processing routine 62 may go on to derive a site traffic matrix (that is, a traffic matrix for all sub-networks) by taking the maximum mean value for each source station/destination station pairing, to be found amongst all the sub-network traffic matrices and using this value as the entry for that pairing in the site traffic matrix.

The "station allocation" routine is called by the main program following processing of the traffic matrices, is illustrated in Figure 9. This routine operates by taking each station in turn and searching through all the sub-network traffic matrices to identify the sub-network carrying the most traffic sourcing from the station concerned; this sub-network is then treated as the one to which the station is directly connected and an appropriate link is made in the linked list linking each sub-network to its associated stations. More particularly, following the start up of the routine (step 80) two variables are initialised to zero, these being the variables "CC-SUM" and "CC-ID" (see step 81). The variable CC-ID is used to hold the identity (SN-ID) of the sub-network which at any particular point during execution of station allocation routine is the current choice for being the sub-network associated with a station under consideration; the variable CC-SUM holds the sum of packets sourcing from the station under consideration that was calculated for the current choice sub-network.

Following initialisation of the variables CC-ID and CC-SUM, the station allocation routine selects in turn each station in the station list 52 and effects steps 83 to 88 for each station. This successive selection and processing is achieved by a loop set up by steps 82 and 89 in combination. For each station (identity "S-ID"), each sub-network contained in the sub-network list 51 is considered in turn and its traffic matrix examined, this successive examination being carried out in the loop established by steps 83 and 87 in cooperation. For each sub-network traffic matrix examined, the traffic sourcing from the currently selected station is added up for all destination stations and the resulting number is stored to variable "SUM" (step 84). The value of "SUM" is then compared with the value held in "CC-SUM" (step 85); if the former is greater than the latter then the value held in "SUM" is transferred to "CC-SUM" and the sub-network identity "SN-ID" associated with the sub-network traffic matrix under consideration is stored in the variable "CC-ID" (step 86).

The comparison of "SUM" and "CC-SUM" is done using hypothesis testing techniques in view of the statistical nature of the quantities concerned. Such techniques are well known to persons skilled in the art and can be found in many standard works (for example, "Introductory Probability and Statistical Applications", Meyer, published by Addison-Wesley, 2nd Edition 1970). Accordingly, a description of the processing involved will not be given here except to note that where the compared mean values are within a certain range of each other, a judgement will be made that no decision can be validly made as to which is the greater. Where such a judgement is made, this is noted since it implies that no valid allocation can be made to associate the station concerned with a particular sub-network (of course, if a subsequent

value of "SUM" for the station does exceed "CC-SUM" sufficiently to enable a statistically valid decision to be made, then this will override any previous 'indeterminable' judgement made for the station).

Step 87 loops the process until all sub-network traffic matrices have been examined in respect of a particular station. Upon exiting of this loop, the variable "CC-ID" will contain the identity of the sub-network whose traffic matrix contained the highest number of entries for traffic sourcing from the station under consideration. This sub-network is taken to be one to which the station under consideration is connected (unless, as indicated, a judgement has been made that a decision cannot be validly made). Therefore, where the station under consideration is the first to be allocated to the sub-network identifed by "CC-ID" (as indicated by the null setting of the pointer S-POINTER in the sub-network entry), the S-POINTER for that sub-network entry is set to point to the entry in the station list 52 corresponding to the station (identity "S-ID") under consideration (step 88). Where the station under consideration is not the first to be allocated to the network identifed by "CC-ID", then the chain of stations associated with the sub-network and pointed to by the relevant S-POINTER is followed to the last station entry in the chain and then the NEXT-S POINTER of that station entry is set to point to the entry for the station under consideration.

Thereafter, step 89 causes the routine to loop back to step 82 until all the stations have been allocated to sub-networks. When this is achieved, the routine terminates (step 19).

After allocating stations to sub-networks, the main program (Figure 7) initiates the "tests for bridges" routine in order to ascertain which pairings of sub-networks are interconnected by bridges. Before describing this routine in detail, the basis for the tests used to identify bridges will be discussed.

Consider the case where it is desired to ascertain whether a bridge directly connects two sub-networks Y and X such as to allow non-local traffic on sub-network Y to pass to sub-network X. If the remaining sub-networks of the network (that is, the group of sub-networks other than X and Y) are considered as a group R, then the network and the flow of traffic from sub-network Y can be represented by the diagram of Figure 10A in the case where a bridge does directly connect sub-networks Y and X, and by the diagram of Figure 10B in the case where there is no bridge connecting sub-networks Y and X. In these diagrams, various traffic flow components are represented by the letter 'T' followed in parenthesis by letters representing the sequence of sub-networks followed by the traffic component from its source on sub-network Y, to its destination. Thus the component of traffic sourcing on sub-network Y and passing to sub-network X through the sub-network group R, is represented by T(YRX).

The dashed traffic flow components T(YRX) and T(YRXR) shown in Figure 10A, will only be present where the network topology includes a loop providing both a direct and an indirect path to sub-network X. Indeed, the traffic flow components T-(YR) and T(YXR) of Figure 10A may also not exist if the sub-networks R are only connected to sub-network Y through sub-network X.

In the first of the three tests for a bridge, consideration is given to what elements of the total traffic coming from sub-network Y (regardless of destination) are seen on the sub-network X and on the sub-network group R. In this consideration, the following representation is used:

T_1 = all traffic from sub-network Y, regardless of destination;

$[T_1]_X$ = that part of T_1 seen on sub-network X;

$T_1(Y..)$ = that component of T_1 following the route indicated in parenthesis.

By reference to Figure 10A, it can be seen that when a bridge directly links sub-network Y to sub-network X:

$$[T_1]_X = T_1(YX) + T_1(YXR) + T_1(YRX) + T_1(YRXR) \quad (3)$$

$$[T_1]_R = T_1(YR) + T_1(YXR) + T_1(YRX) + T_1(YRXR) \quad (4)$$

Formulae (3) and (4) represent the most general cases. The last two components of both formula (3) and (4) will not be present in a non-looped network where there is only a direct path to sub-network X. Furthermore, these components, and also the component $T_1(YR)$ from formula (4) will be absent if the only bridge from sub-network Y is the one connecting to sub-network X.

Referring now to Figure 10B, it can be seen that in the absence of a bridge directly linking sub-network Y to sub-network X:

$$[T_1]_X = T_1(YRX) + T_1(YRXR) \quad (5)$$

$$T_1]_R = T_1(YRX) + T_1(YRXR) + T_1(YR) \quad (6)$$

from which it follows that, absent a bridge connecting sub-network Y to sub-network X,

$$[T_1]_R >= [T_1]_X \quad (7)$$

If formula (5) is not satisfied, ie. if:

$[T_1]_X > [T_1]_R$ (8)

then a bridge must exist directly connecting sub-network Y to sub-network X. Stated in words, a bridge can be taken to exist between sub-network Y and sub-network X, if there is more traffic from stations on sub-network Y seen on sub-network X, than seen on all the sub-networks of the remaining group R. This test will not necessarily identify a bridge which exists between sub-networks Y and X, but it will do so in many cases and will not find a bridge where none exists.

The second of the three tests for a bridge is similar to the first but only considers the traffic sourcing from sub-network Y that is not destined for sub-network X, this traffic being designated T_2. By reference to the traffic components shown in Figure 10B not destined for sub-network X, it can be seen that in the absence of a brdige linking sub-network Y to sub-network X:

$[T_2]_X = T_2(YRXR)$ (9)

$[T_2]_R = T_2(YR) + T_2(YRXR)$ (10)

From this it follows that in the absence of a bridge:

$[T_2]_R >= [T_2]_X$ (11)

if formula (11) is not satisfied, i.e. if:

$[T_2]_X > [T_2]_R$ (12)

then a bridge must exist directly connecting sub-network Y to sub-network X. In other words, a bridge can be taken to exist between sub-network Y and sub-network X, if there is more traffic from stations on sub-network Y not destined for sub-network X, seen as sub-network X, than seen on all the sub-networks of the remaining group R. Again, this test will not necessarily find a bridge but it will not find one if one does not exist.

The third of the three tests for the presence of a bridge relies on the fact that most practical bridges are 'leaky', that is, they allow a small portion of the traffic destined for a sub-network connected to the bridge to cross through the bridge to a connected sub-network. This leaked traffic was not taken into account in considering the traffic components relevant to the first two tests as the leaked components will be small in comparison with the other components and, in any case, as will be apparent from what follows, the presence of these leaked components does not alter the validity of the previously derived test formulae.

If the traffic sourcing from sub-network Y and destined from sub-network Y is designated T_3, then the possible traffic components are as shown in Figure 11. It should be noted that the fate of the 'leaked' traffic is uncertain and will depend on the nature of the operation of the bridges of the network. Thus, whilst the bridges may serve to filter out the leaked traffic from further transmission through the network, this may not always be the case. Accordingly, Figure 11 includes all cases of the leaked traffic being allowed to return to sub-network Y as well as all cases of the leaked traffic being filtered out.

From Figure 11 it can be seen that in the absence of a bridge linking sub-network Y to sub-network X:

$[T_3]_X = T_3(YRXRY) + T_3(YRX) + T_3(YRXR)$ - (13)

$[T_3]_R = T_3(YRXRY) + T_3(YRX) + T_3(YRXR) + + T_3(YR) + T_3(YRY)$ (14)

From this it follows that in the absence of a bridge:

$[T_3]_R >= [T_3]_X$ (15)

If formula (15) is not satisfied, ie. if:

$[T_3]_X > [T_3]_R$ (16)

then a bridge must exist directly connecting sub-network Y to sub-network X. In other words, a bridge can be taken to exist between sub-network Y and sub-network X, if there is more traffic from stations on sub-network Y destined for sub-network Y, seen on sub-network X, than seen on all the sub-networks of the remaining group R. Again, this test will not necessarily find a bridge but it will not fine one if one does not exist.

It will be appreciated that in view of the direction of the inequality in formula (16), superimposing the leakage traffic flows onto those discussed above in respect of the first and second bridge tests, does not alter the results obtained in formulas (8) and (12).

The "test for bridges" routine shown in Figure 12 takes each pairing of sub-networks and then applies each of the three tests for bridges described above to ascertain whether or not a bridge is present linking a first one of the sub-networks to the second. The nature of the tests is such that they are uni-directional - that is, they test for whether there is a bridge connection allowing traffic flow in one direction from one sub-network to the other. However, the routine of Figure 12 applies the tests in both directions by first applying the tests to the sub-network pair in one ordering and then subsequently applying the tests to the sub-network pair in the ordering.

Considering Figure 12 in more detail, each sub-network of the sub-network list 51 is taken in turn as the first sub-network Y of a sub-network pair to be subjected to the bridge tests (see step 91). For each sub-network assigned as the sub-network Y, each of the remaining sub-networks is taken in turn as the second sub-network X of the pair of sub-networks to be tested (step 93). Once the pair of sub-networks Y,X to be tested has been set up, the first test for a bridge is applied by calling an appropriate routine (step 94). After this first test has been completed, a check is made to see whether the test were successful, if it was, then the second and third tests are not applied to the pair Y,X. However, if a bridge was not found by the first test, the second test is then executed (step 96).

If after completion of the second test, a bridge still has not been found between pairing YX (step 97), then the third test for the presence of the bridge is carried out (step 98).

The situation is then reached when either a bridge has been found between the sub-network pair YX, or all three tests have been carried out and no bridge has been found. Where a bridge has been found, the bridge list will have been updated as the final step of the relevant bridge test sub-network.

Thereafter, the next sub-network pairing is taken, either by looping back to select the next second sub-network X (step 99), or where all the sub-networks (except for the current Y sub-network) have been used as the second sub-network X, by looping back to select a new sub-network for the first sub-network Y (step 100).

In due course, all sub-network pairings will have been tested in both directions. Thereafter, the "test for bridges" routine is terminated.

Figure 13 illustrates the "first test for bridge" sub routine. It will be recalled that this first test involves comparing the total of all the traffic sourcing from a first sub-network Y as seen on a second sub-network X, with the total of all the traffic sourcing from the sub-network Y as seen on all the other sub-networks (that is, excepting X and Y). The first step 102 of the sub routine involves accessing the traffic matrix for the sub-network X and summing all traffic sourcing from sub-network Y (this traffic being the traffic originating from the stations on sub-network Y, these stations being identified by reference to the linked list starting at the S-POINTER for sub-network Y). The resultant sum is stored to a variable representing $[T_1]_X$.

Next, in step 103 and 104 the traffic matrix for each sub-network other than the networks Y and X is examined in turn and an accumulative total arrived for traffic sourcing from sub-network Y, no regard being paid to the fact that the same traffic sourcing from sub-network Y may be viewed on several of the sub-networks under consideration. The total traffic value derived from all these sub-network traffic matrices is stored in a variable representing $[T_1]_R$. Next, the values of $[T_1]_X$ and $[T_1]_R$ are compared (step 105); in the present case, in view of the statistical nature of the data being processed, this comparison is done using hypothesis testing techniques. If the value of $[T_1]_X$ is greater than the value of $[T_1]_R$ then a decision is made that a bridge exists permitting traffic flow directly from sub-network Y to sub-network X; a corresponding entry is therefore made in the bridge list identifying the new bridge and the first and second sub-networks that it links (step 106). Thereafter, "the first test for bridge" sub-routine terminates.

The "second test for bridge" sub-routine is similar to that shown in Figure 13 except that only traffic sourcing from sub-network Y that is not destined for sub-network X is totalled in the steps corresponding to steps 102 and 103.

The "third test for a bridge" sub-routine is also similar to that shown in Figure 13 but with only traffic sourcing from sub-network Y that is also destined for sub-network Y being totalled in the steps corresponding to the steps 102 and 103.

Many variations can be made to the above-described topology discovery method. For example, the or each test for a bridge can be applied in respect of traffic from any particular station on the sub-network Y of a sub-network pair Y,X under consideration and the test can be repeated for traffic from each station on the sub-network Y. Furthermore, the traffic matrix data may be derived by the monitoring devices themselves and then passed to the processing station 13 at the end of the sampling period; in this case, the monitoring devices could record each and every packet rather than sampling the packets. Where only particular topology features are required or only particular tests are to be used, the provision of the full range of traffic information such as is contained in the sub-network traffic matrices is unnecessary. For example, where only the first test for a bridge is to be used and the source information contained in a packet indicates directly the originating sub-network, the monitoring devices can directly record how many packets are observed sourcing from a particular sub-network and this data can then be supplied to the station 13 for immediate use in the "first test for a bridge" sub-routine (Figure 13).

Although the above disclosed topology discovery method has been described in relation to an Ethernet-type network, it will be appreciated that it can also be applied to other networks, such as an Internet network, to ascertain network topology features.

The foregoing topology discovery method has been described in terms of data structures that can be readily implemented using traditional programming techniques that consider programs and data as separate entities; however, it will be appreciated that the described method can be implemented using techniques such as object-oriented programming.

Claims

1. A method of ascertaining topology features of a network that comprises a plurality of sub-networks, spanning devices (14) interconnecting the sub-networks, and stations (11) operative to source and/or sink traffic to/from the sub-networks, the traffic being in the form of discrete message packets each including source and destination information (15), characterised in that said method comprises the steps of:
   - monitoring the traffic on at least some of said sub-networks to derive data concerning at least the amounts and origins of the traffic so monitored, and
   - processing the said data to determine selected topology features of the network.

2. A method, according to claim 1, of ascertaining to which said sub-network a selected said station (11) is connected where said source information included in said message packets serves to identify the stations sourcing the packets, said monitoring step (40-47, 65-75) serving to determine the amount of traffic originating from said selected station (11) that is carried by each said sub-network, and said processing step (83-87) involving determining the sub-network carrying the most traffic from said selected station (11), this sub-network being taken as the one to which said selected station (11), is connected.

3. A method, according to claim 1, of ascertaining to which said sub-networks said stations (11) are connected where said source information included in said message packets serves to identify the stations sourcing the packets, said method involving carrying out the method of claim 2 for each said station (11) with the monitoring step (65-75) being carried out for all stations concurrently.

4. A method, according to claim 1, of ascertaining whether a first said sub-network (Y) is directly connected by a said spanning device (14) to a second said sub-network (X), wherein:
   - said method includes the step of providing association information (40-47, 65-75, 63) enabling the said source information included in a message packet to be used to identify a message packet as having originated from said first sub-network (Y);
   - said monitoring step (40-47, 65-75, 63) involves monitoring the traffic on said second sub-network (X) and on each sub-network of the group (R) of sub-networks made up of all said sub-networks other than said first and second sub-networks, whereby to derive, with the use of said association information, data on the amounts of traffic sourcing from said first sub-network (Y) that are carried by said second sub-network (X) and the sub-networks of said group (R) of sub-networks; and
   - said processing step (94) is operative to use said data to compare the said amounts of traffic ($[T_1]_X$,$[T_1]_R$) sourcing from said first sub-network (Y), that are respectively carried by said second sub-network (X) and by the sub-networks of said group (R) taken together, the processing step determining that a spanning device (14) directly connects the first sub-network (Y) to the second (X) when the larger of these amounts of traffic is carried by said second sub-network (X).

5. A method, according to claim 1, of ascertaining whether a first said sub-network (Y) is directly connected by a said spanning device (14) to a second said sub-network (X), wherein:
   - said method includes the step of providing association information (40-47, 65-75, 63) enabling:
     (a) the said source information included in a message packet to be used to identify a message packet as having originated from said first sub-network (Y); and
     (b) the said destination information included in a message packet to be used to identify a message packet as having a destination other than said second sub-network (X);
   - said monitoring step (40-47, 65-75, 63) involves monitoring the traffic on said second sub-network (X) and on each sub-network of the group (R) of sub-networks made up of all said sub-networks other than said first and second sub-networks, whereby to derive, with the use of said association information, data on the amounts of traffic sourcing from said first sub-network (Y) and having a

destination other than on said second sub-network (X), that are carried by said second sub-network (X) and the sub-networks of said group (R) of sub-networks; and

- said processing step [96] is operative to use said data to compare the said amounts of traffic ($[T_2]_X$,$[T_2]_R$) sourcing from said first sub-network (Y) and having a destination other than on said second sub-network (X), that are respectively carried by said second sub-network (X) and by the sub-networks of said group (R) taken together, the processing step determining that a spanning device (14) directly connects the first sub-network (Y) to the second (X) when the larger of these amounts of traffic is carried by said second sub-network (X).

6. A method, according to claim 1, of ascertaining whether a first said sub-network (Y) is directly connected by a said spanning device (14) to a second said sub-network (X) where the spanning devices (14) used in the network do not act perfectly to block messages packets with a destination on a connected sub-network from passing through the spanning device (14) from that sub-network to another, wherein:

- said method includes the step of providing association information (40-47, 65-75, 63) enabling:
  (a) the said source information included in a message packet to be used to identify a message packet as having originated from said first sub-network (Y); and
  (b) the said destination information included in a message packet to be used to identify a message packet as having a destination on said first sub-network (Y);
- said monitoring step (40-47, 65-75, 63) involves monitoring the traffic on said second sub-network (X) and on each sub-network of the group (R) of sub-networks made up of all said sub-networks other than said first and second sub-networks, whereby to derive, with the use of said association information, data on the amounts of traffic sourcing from said first sub-network (Y) and having a destination on said first sub-network (Y), that are carried by said second sub-network (X) and the sub-networks of said group (R) of sub-networks; and
- said processing step (98) is operative to use said data to compare the said amounts of traffic ($[T_3]_X$,$[T_3]_R$) sourcing from said first sub-network (Y) and having a destination on said first sub-network (Y), that are respectively carried by said second sub-network (X) and by the sub-networks of said group (R) taken together, the processing step determining that a spanning device (14) directly connects the first sub-network (Y) to the second (X) when the larger of these amounts of traffic is carried by said second sub-network (X).

7. A method according to claim 4, 5 or 6, wherein said processing step (94; 96; 98) derives the relevant amount of traffic sourcing from said first sub-network (Y) that is carried by the sub-networks of said group (R) taken together, by adding together the corresponding traffic amounts for each sub-network of said group (R) without regard to any duplication that may occur as a result of the same message packet being monitored on more than one sub-network of the group.

8. A method according to claim 4, wherein said association information further enables the said destination information included in a message packet to be used to identify the destination sub-network of the packet, said processing step (94) deriving the said amount of traffic sourcing from said first sub-network (Y) that is carried by the sub-networks of said group (R) taken together, by:

- determining for each sub-network of said group (R), the components of the sub-network traffic sourcing from said first sub-network (Y), that are destined for that sub-network and for said second sub-network (Y), and
- adding together the sub-network traffic components so determined for all the sub-networks of said group (R).

9. A method according to claim 5 or 6, wherein said processing step derives the relevant said amount of traffic sourcing from said first sub-network (Y) that is carried by the sub-networks of said group (R) taken together, by:

- determining for each sub-network of said group (R), the components of the relevant sub-network traffic sourcing from said first sub-network (Y), that are destined for that sub-network and for said second sub-network (X), and
- adding together the relevant sub-network traffic components so determined for all the sub-networks of said group (R).

10. A method according to claim 4, 5 or 6, wherein

said step of providing association information (40-47, 65-75, 63) involves the method of claim 2 or 3 for associating stations with said first sub-network, the monitoring step effected by the method of claim 2 or 3 being carried out either independently of, or in conjunction with, the monitoring step recited in the relevant one of claims 4, 5 or 6.

11. A method according to claim 4, 5 or 6, wherein said association information further enables the said source information included in a message packet to be used to identify from which station on said first sub-network a said message packet has originated; the said monitoring and processing steps for determining whether a spanning device (14) is present between said first and second sub-networks (Y,X), being carried out in respect of traffic from the or each station of said first sub-network taken separately.

12. A method of ascertaining whether two sub-networks of a network are directly connected by a bi-directional spanning device (14), said method involving carrying out the method of claim 4, 5 or 6 in respect of traffic sourcing from one said sub-network and then repeating that method for traffic sourcing from the other said sub-network.

13. A method of ascertaining the topology of at least a portion of a network in terms of its sub-networks and their interconnection by spanning devices (14), said method involving carrying out the method of any one of claims 4, 5 or 6, in respect of each pairing of sub-networks of the network or the portion thereof under consideration, whereby to determine which of the sub-networks are interconnected by spanning devices.

14. A method according to any one of the preceding claims, wherein said monitoring step (40-47, 65-75, 63) involves the generation of traffic matrices for each sub-network.

15. A method according to any one of the preceding claims, wherein said monitoring step (40-47, 65-75, 63) involves sampling the traffic on the sub-networks, said processing step (94; 96; 98) effecting its determination by using hypothesis testing.

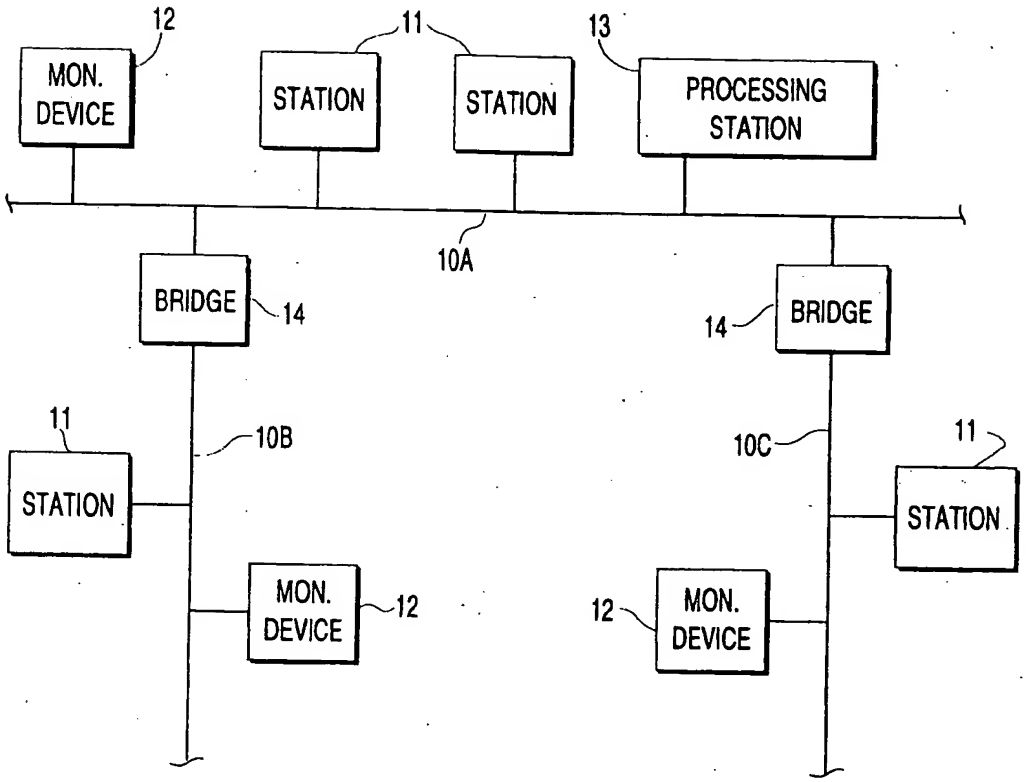

12
MON. DEVICE
11
STATION
STATION
13
PROCESSING STATION
10A
BRIDGE
14
14
BRIDGE
11
STATION
10B
10C
11
STATION
MON. DEVICE
12
12
MON. DEVICE

FIG 1

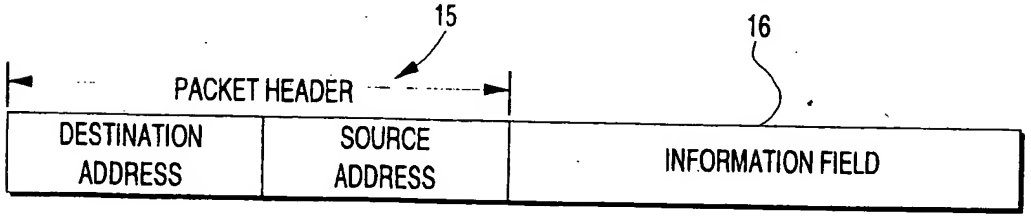

15
16
PACKET HEADER
DESTINATION ADDRESS
SOURCE ADDRESS
INFORMATION FIELD

FIG 2

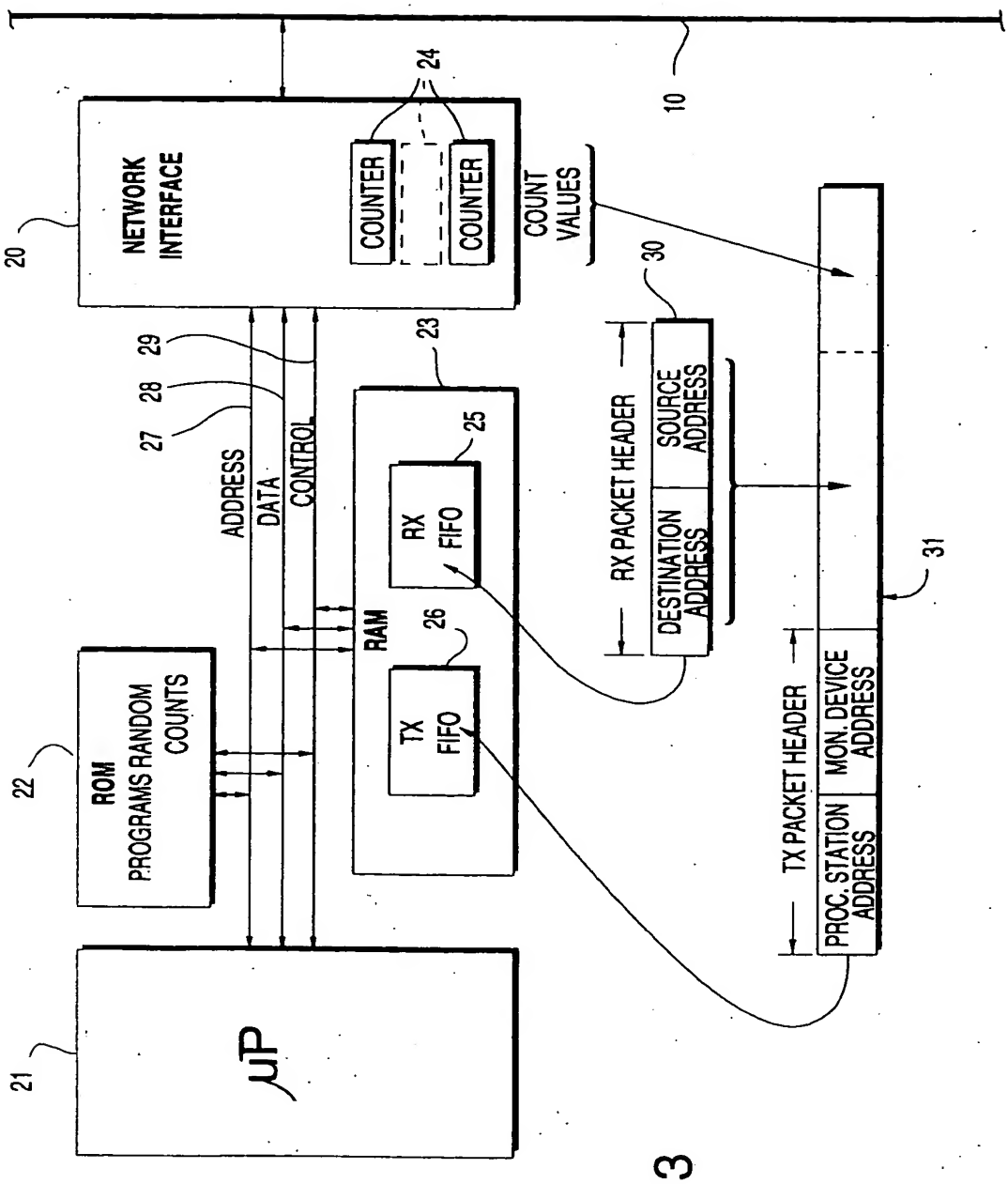

10
20
NETWORK INTERFACE
24
COUNTER
COUNTER
COUNT VALUES
30
RX PACKET HEADER
SOURCE ADDRESS
DESTINATION ADDRESS
31
TX PACKET HEADER
MON. DEVICE ADDRESS
PROC. STATION ADDRESS
23
RAM
25
RX FIFO
26
TX FIFO
29
28
27
ADDRESS
DATA
CONTROL
22
ROM
PROGRAMS RANDOM COUNTS
21
μP

FIG 3

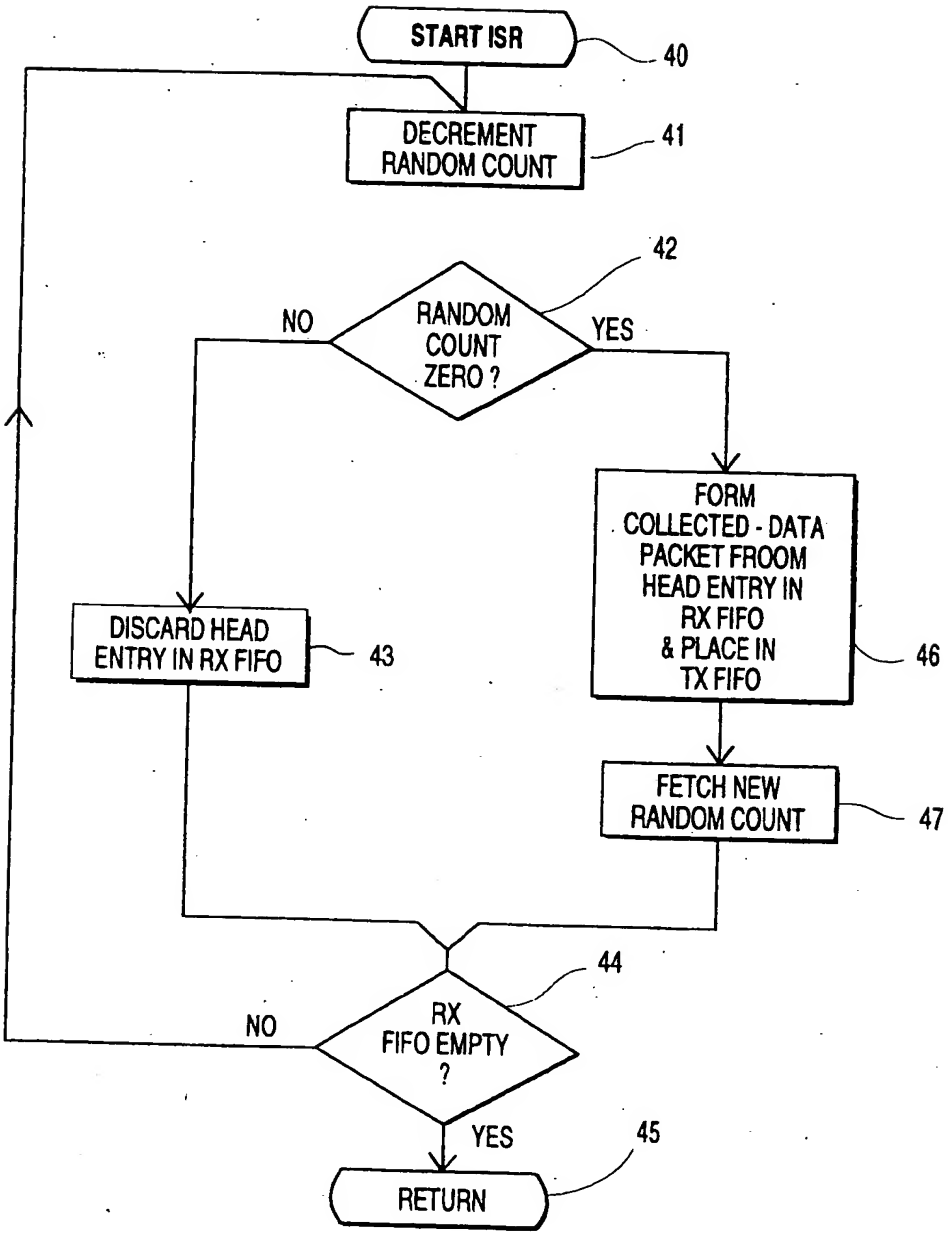

START ISR
40
DECREMENT RANDOM COUNT
41
42
RANDOM COUNT ZERO ?
NO
YES
DISCARD HEAD ENTRY IN RX FIFO
43
FORM COLLECTED - DATA PACKET FROOM HEAD ENTRY IN RX FIFO & PLACE IN TX FIFO
46
FETCH NEW RANDOM COUNT
47
44
RX FIFO EMPTY ?
NO
YES
45
RETURN

FIG 4

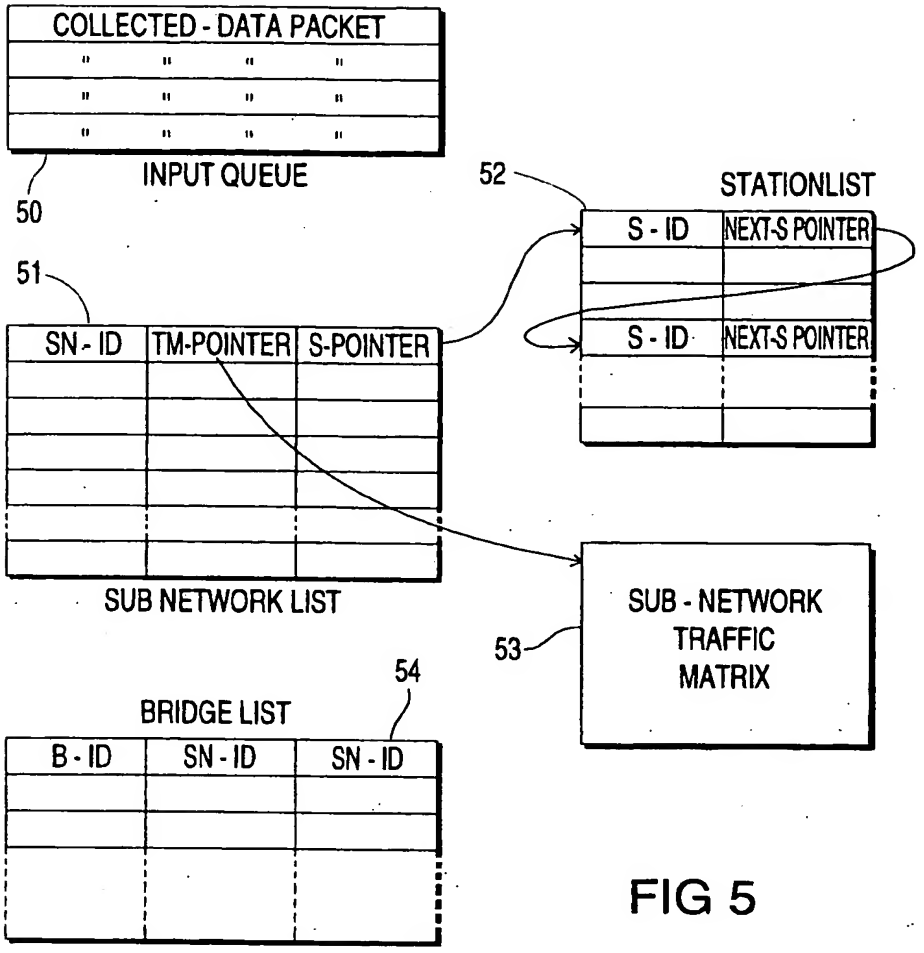
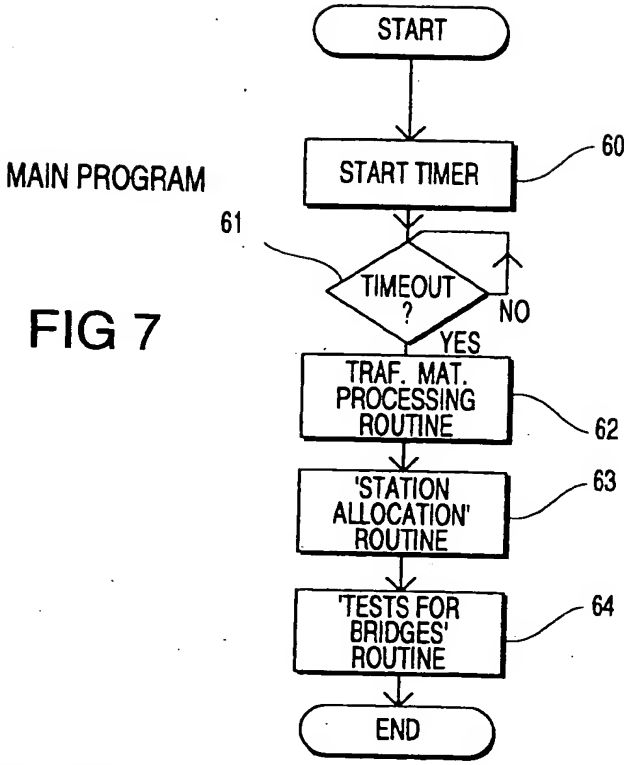


FIG 5

FIG 6

SUB - NETWORK TRAFFIC MATRIX		DESTINATION STATION			
		11A	11B	11C - - - - - -	11N
	11A	–	21	9	65
SOURCE	11B	42	–	100	100
STATION	11C	69	84	–	15
	⋮				
	11N	150	29	75	–



MAIN PROGRAM

FIG 7

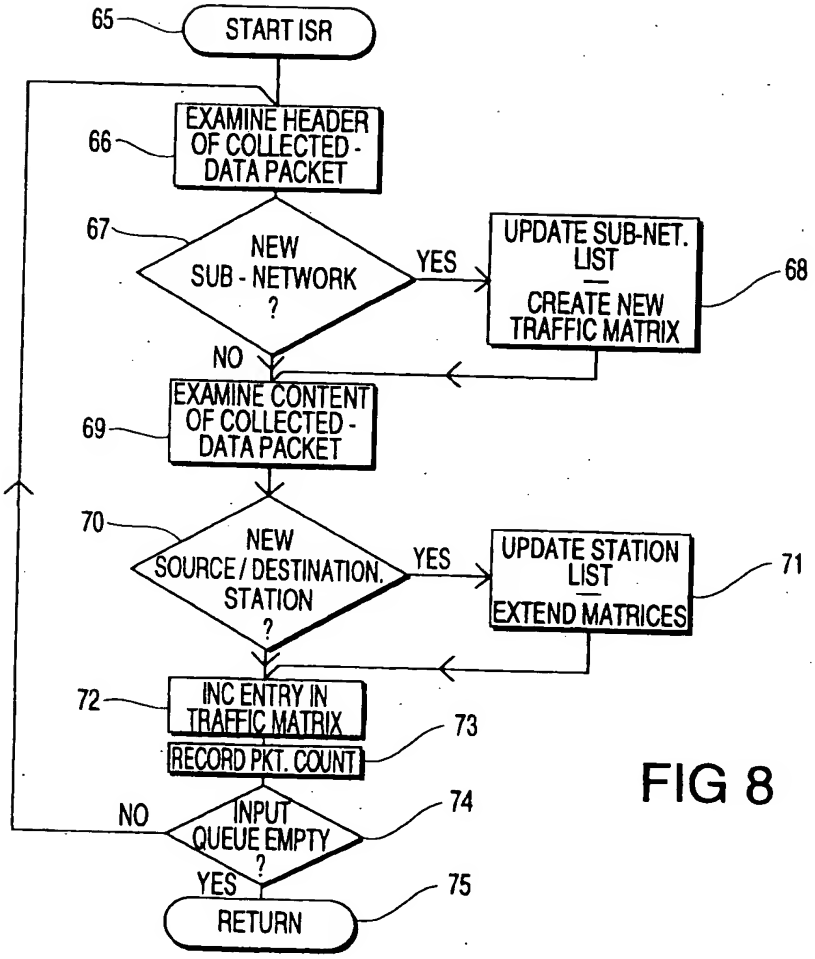
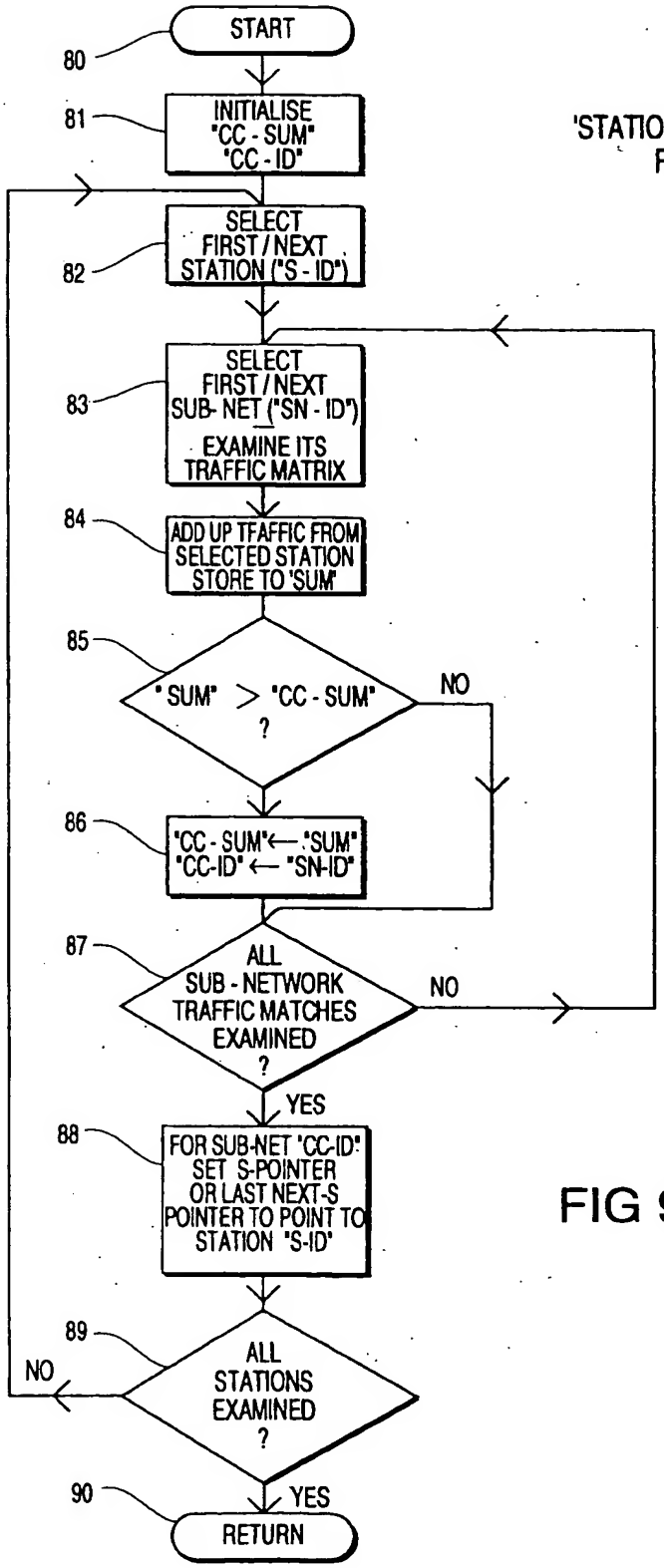


FIG 8



'STATION ALLOCATION' ROUTINE

FIG 9

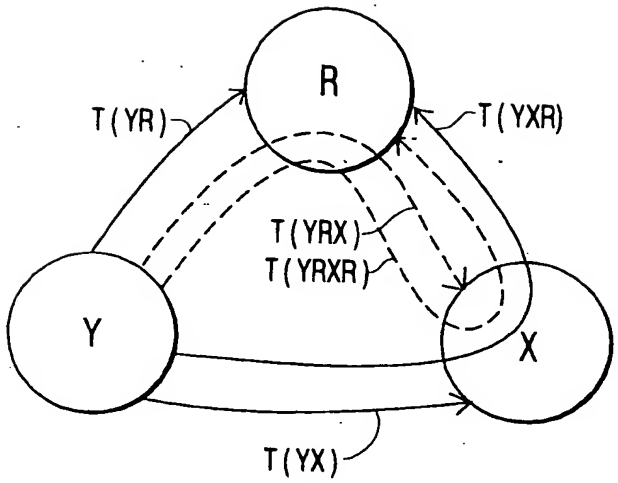


FIG 10A

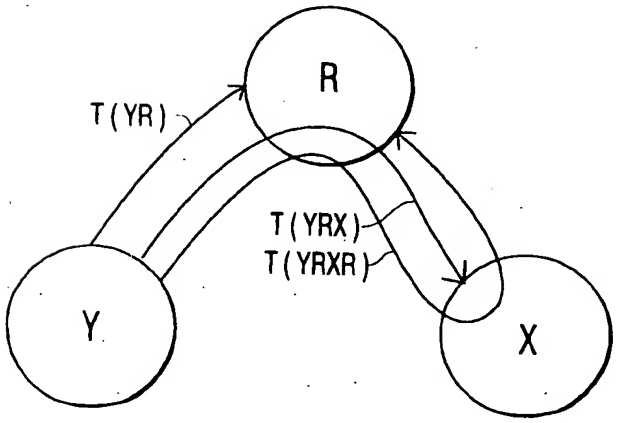


FIG 10B

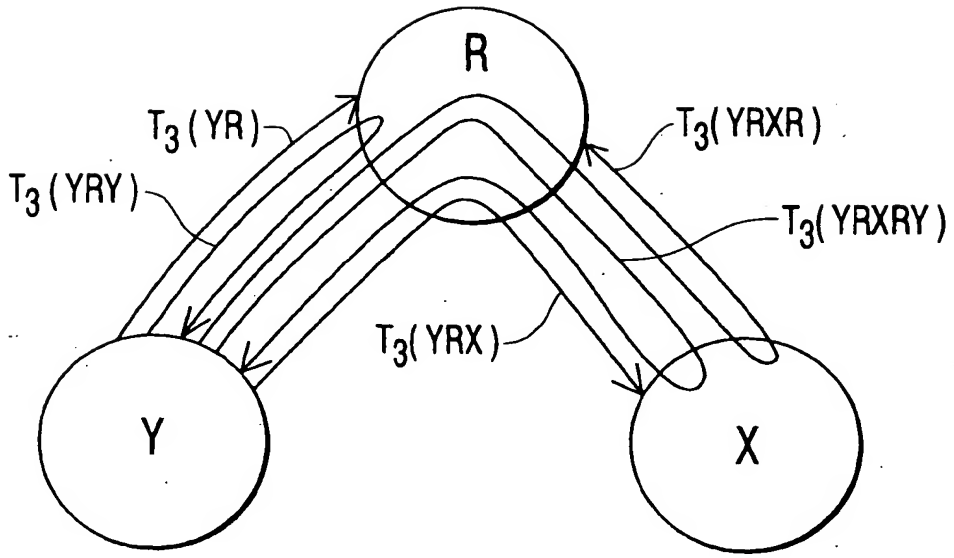
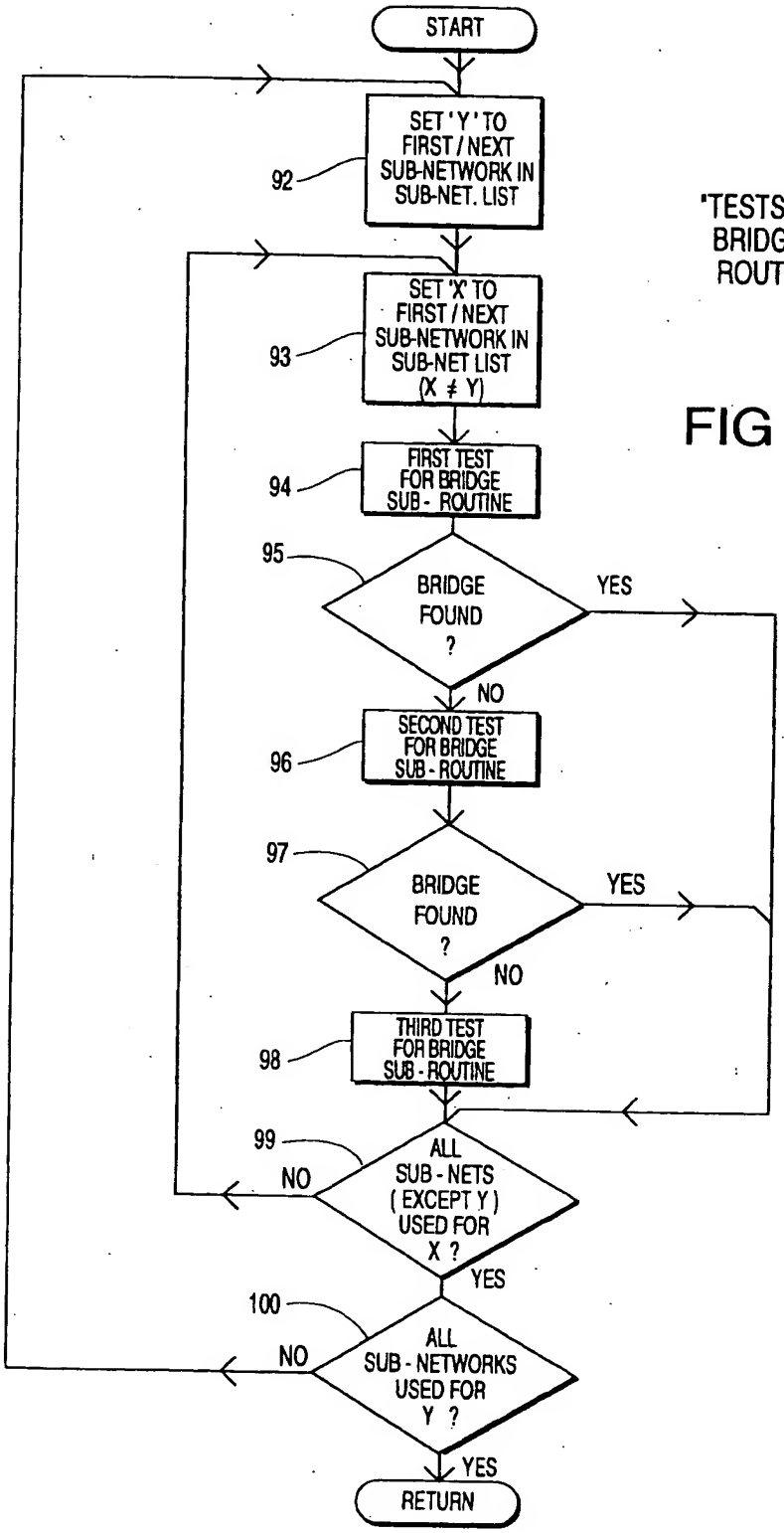
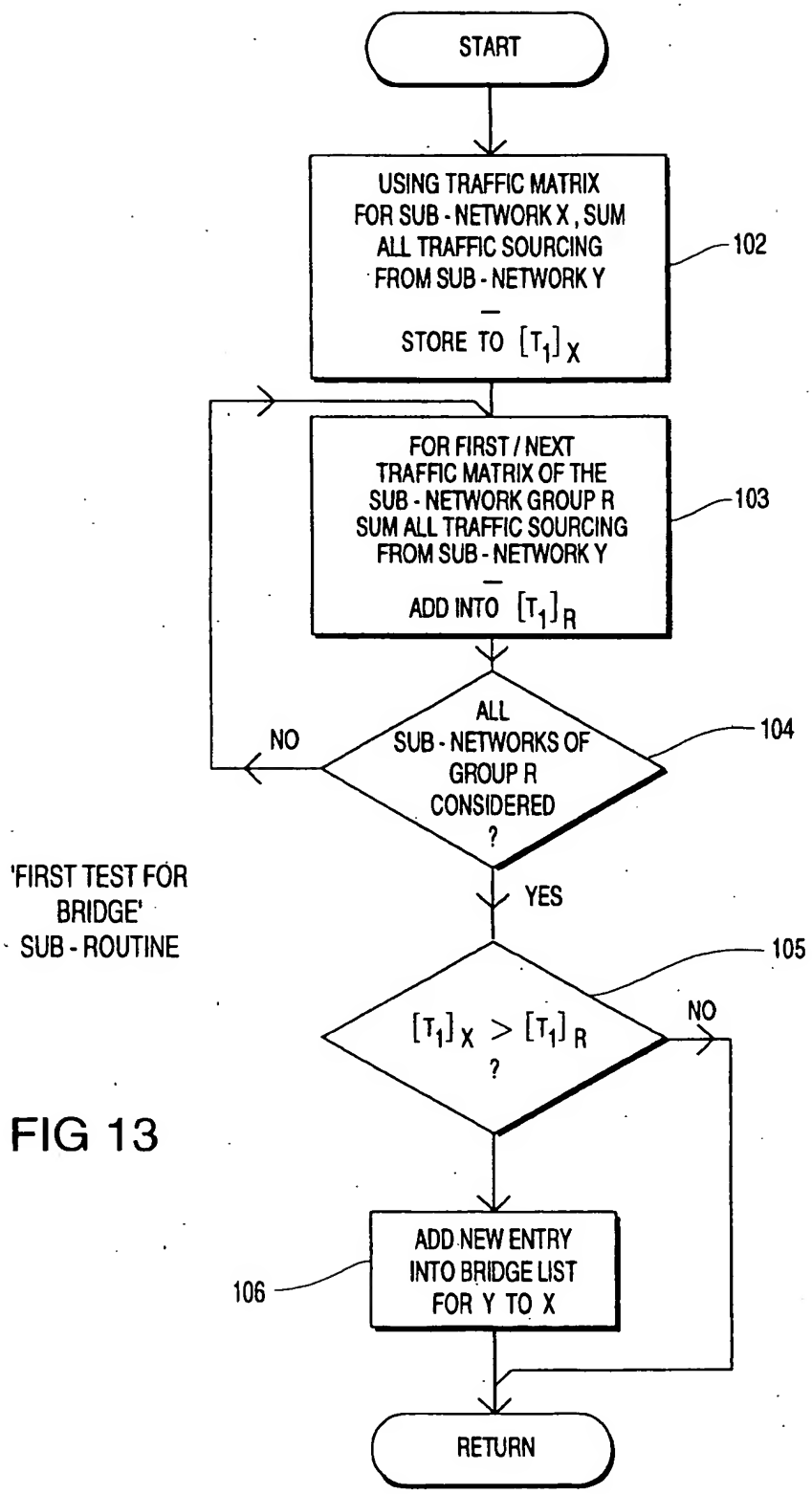


FIG 11



'TESTS FOR BRIDGES' ROUTINE

FIG 12



'FIRST TEST FOR BRIDGE' SUB - ROUTINE

FIG 13

European Patent Office

EUROPEAN SEARCH REPORT

Application Number

EP 91 30 1204

DOCUMENTS CONSIDERED TO BE RELEVANT

Category	Citation of document with indication, where appropriate, of relevant passages	Relevant to claim	CLASSIFICATION OF THE APPLICATION (Int. Cl.5)
A	FUTURE GENERATIONS COMPUTER SYSTEMS. vol. 5, no. 4, January 1990, AMSTERDAM NL pages 359 - 363; U. DE CARLINI ET AL: 'THE MONITORING OF INTER-PROCESS COMMUNICATION IN DISTRIBUTED SYS-TEMS' * paragraph 1 3 * * figures 3,4 * *	1	H 04 L 12/26
A	EP-A-0 265 106 (HEWLETT PACKARD) * column 2, line 41 - column 3, line 40 * * column 5, line 33 - column 6, line 24 * * column 8, line 28 - line 39 * * figures 2,4,5 * *	1	
A	IEEE NETWORK: THE MAGAZINE OF COMPUTER COM-MUNICATIONS. vol. 1, no. 3, July 1987, NEW YORK US pages 32 - 39; D.RITTER ET AL: 'A MULTI-PURPOSE, DISTRIBUTED LAN TRAFFIC MONITORING TOOL' * the whole document * *	1	
A	IEEE NETWORK: THE MAGAZINE OF COMPUTER COM-MUNICATIONS. vol. 1, no. 3, July 1987, NEW YORK US pages 13 - 19; M. SOHA: 'A DISTRIBUTED APPROACH TO LAN MONITORING USING INTELLIGENT HIGH PERFOR-MANCE MONITORS'	1	TECHNICAL FIELDS SEARCHED (Int. Cl.5) H 04 L

The present search report has been drawn up for all claims

Place of search	Date of completion of search	Examiner
The Hague	30 January 92	CANOSA ARESTE C.

CATEGORY OF CITED DOCUMENTS

X: particularly relevant if taken alone
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E: earlier patent document, but published on, or after the filing date
D: document cited in the application
L: document cited for other reasons

&: member of the same patent family, corresponding document